

The Jellyfish Algorithm

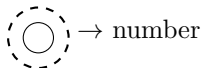
Ellie Grano

UC Santa Barbara

November 20, 2010

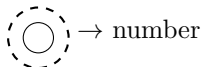
Background

- ▶ Evaluation algorithms common in topology
 - ▶ Examples: Kauffman bracket, HOMFLY polynomial
 - ▶ Idea:



Background

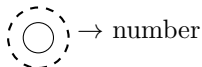
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- ▶ 2008 - Morrison, Peters, and Snyder
Skein theory for the D_{2n} planar algebra

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 - ▶ Idea:



- ▶ 2008 - Morrison, Peters, and Snyder
Skein theory for the D_{2n} planar algebra
- ▶ 2009 - Bigelow
Skein theory for the ADE planar algebras
"jellyfish algorithm" introduced

The Temperley-Lieb planar algebra

- ▶ For each k , \mathcal{TL}_{2k} is an algebra over $\mathbb{C}(q)$. As a vector space, \mathcal{TL}_{2k} is spanned by diagrams with k nonintersecting strands. The multiplication operation is vertical stacking. We also have the following "bubble bursting" relation:

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$$\text{bubble} = \delta \cdot \text{empty circle}, \text{ where } \delta = q + q^{-1}$$

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$$\left[\begin{array}{c} \text{bubble} \\ \text{bubble} \end{array} \right] = \left[\begin{array}{c} \text{cup} \\ \text{circle} \\ \text{cap} \end{array} \right] = \delta \cdot \left[\begin{array}{c} \text{bubble} \\ \text{bubble} \end{array} \right] \in \mathcal{TL}_6$$

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- ▶ These vector spaces assemble together into a planar algebra with $\mathcal{TL}_0 \cong \mathbb{C}(q)$.

The Temperley-Lieb planar algebra

► $\left(\begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array} \right) := iq^{\frac{1}{2}} \left(\begin{array}{c} \diagdown \quad \diagup \\ \diagup \quad \diagdown \end{array} \right) \in \mathcal{TL}_4$

This satisfies R2 and R3.

For R1, we get a positive twist factor $iq^{3/2}$.

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
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- ▶ Jones-Wenzl projections.

For each k there is a unique element $p_k \in \mathcal{TL}_{2k}$ such that:

- ▶ $p_k^2 = p_k$
- ▶ p_k is uncappable.

 = zero

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- ▶ p_k is uncappable.

$$\left(\begin{array}{c} | \quad | \quad | \\ \cap \\ \boxed{p_k} \\ | \quad | \quad | \end{array} \right) = \text{zero}$$

It follows that:

$$p_k = \left| \begin{array}{c} | \quad | \quad | \\ \dots \\ | \quad | \quad | \end{array} \right| + \sum \alpha_Q \cdot Q, \text{ where each } Q \text{ contains a cap.}$$

Temperley-Lieb when q is a root of unity

If $q = e^{i\pi/n+1}$, then p_n becomes negligible. So for \mathcal{TL} at this value of q , we must add the relation $p_n = zero$ (this gives us the A_n planar algebra). For example, if $q = e^{i\pi/6}$, then \mathcal{TL} will have the relations:

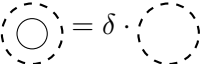

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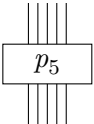
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▶  = $\delta \cdot$ 

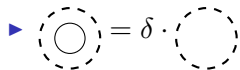
▶  = *zero*

The D_{2n} planar algebra (\mathcal{P})

Fix $n = 2$ and $q = e^{i\pi/6}$. Define \mathcal{P} to be the planar algebra generated by a single S -box in \mathcal{P}_4 subject to the following relations:

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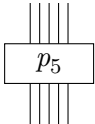
► 

$$\text{Dashed circle with inner solid circle} = \delta \cdot \text{Dashed circle}$$

The D_{2n} planar algebra (\mathcal{P})

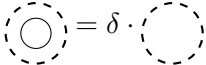
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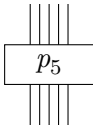
▶  = $\delta \cdot$ 


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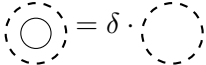

▶  $\delta \cdot$

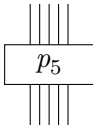
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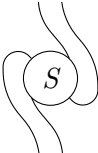

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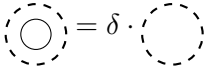

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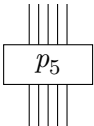
▶  $= zero$


▶  $= i \cdot$ 

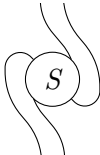

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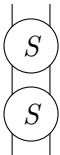
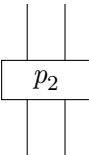
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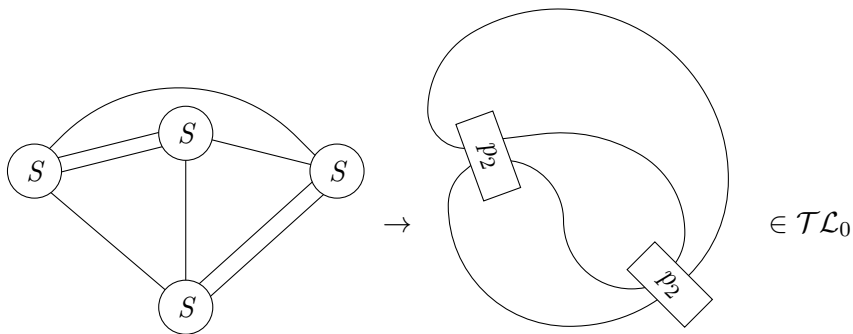
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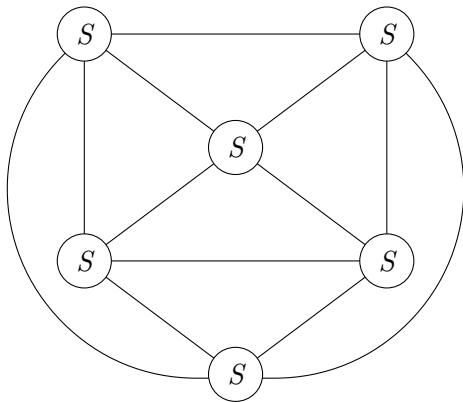
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▶  $=$ 

An example of evaluating a diagram using the relations:



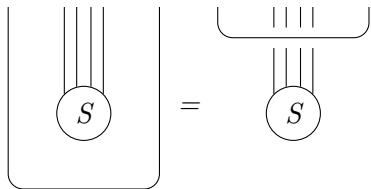
But what about:



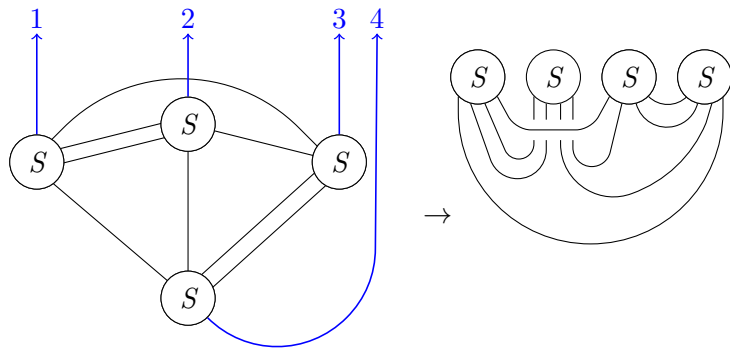
Partial braiding

Theorem

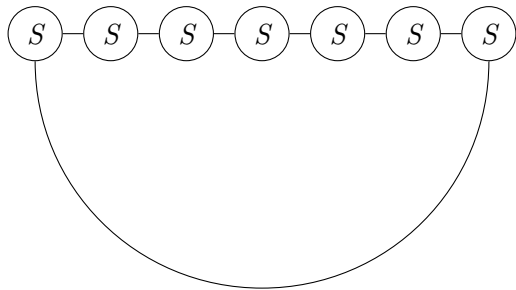
The relations imply the following partial braiding:



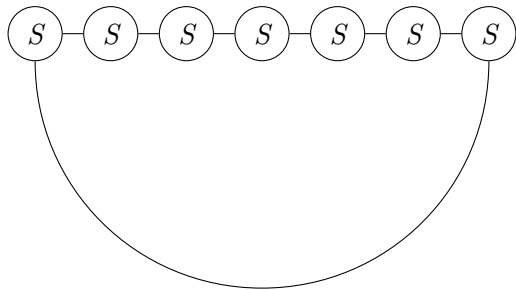
Example of the jellyfish algorithm



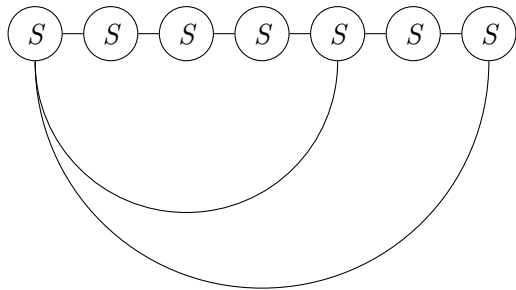
Will there always be an S^2 ?



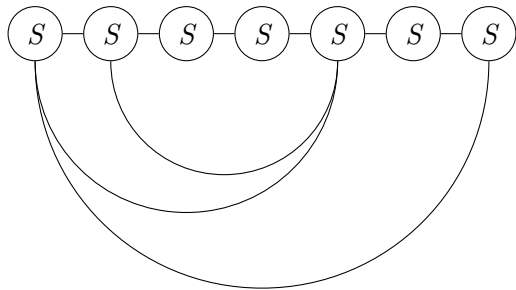
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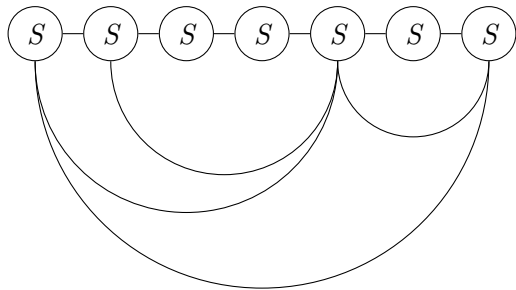
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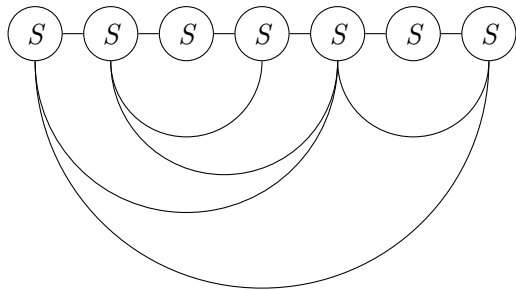
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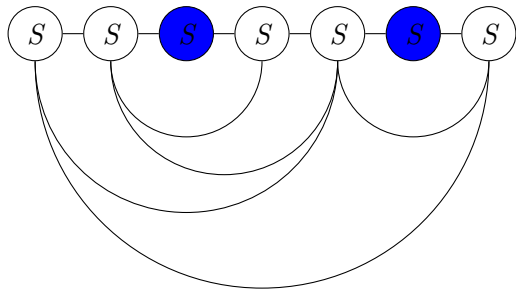
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The jellyfish algorithm

Part I

- ▶ Draw an arc for each S-box
- ▶ Order the arcs
- ▶ Drag the S-boxes in order under any strands
- ▶ Evaluate crossings
- ▶ Go to Part II

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Part II

- ▶ If there are zero S-boxes, evaluate as in \mathcal{TL}
- ▶ If there is a cap on an S-box, evaluate as zero
- ▶ If there are two or more S-boxes
 - ▶ Pick a pair of S-boxes connected by at least two strands
 - ▶ Choose two strands connecting the pair and replace with a p_2
 - ▶ Put in the correct coefficient
 - ▶ Start Part II again

The jellyfish algorithm

Part I

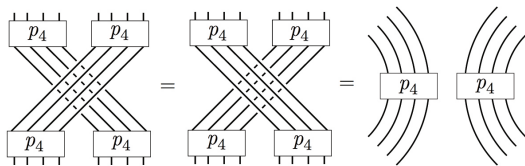
- ▶ **Draw** an arc for each S-box
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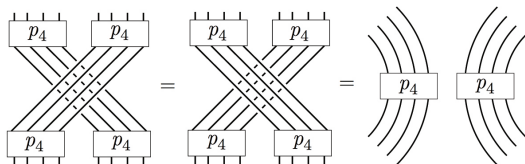
Change the ordering of the arcs

Fact (for our value of q):



Change the ordering of the arcs

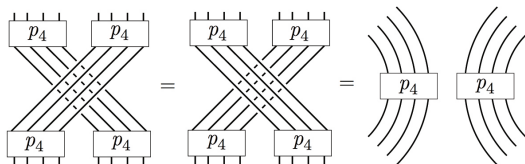
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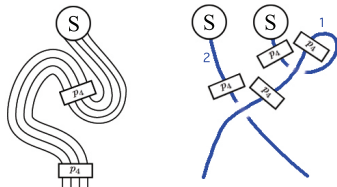
Idea: Add a p_4 along an S -box arc. Do this for each arc crossing that needs to be changed, then use the above fact.

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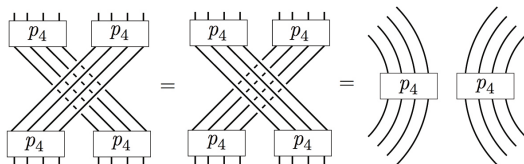


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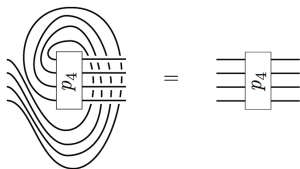


Alter the path of an arc

Notice that



implies:



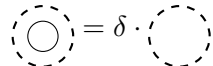
Thus we have all three Reidemeister moves for the S -box arcs.

Check the algorithm respects the relations

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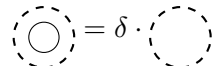
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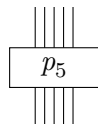
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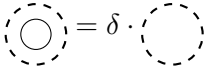

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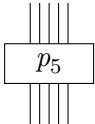
▶  $\delta \cdot$


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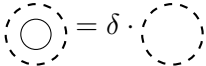

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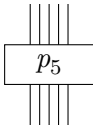
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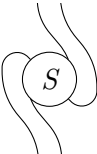

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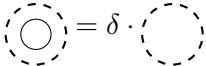

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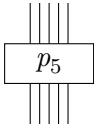
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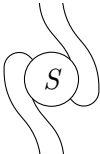

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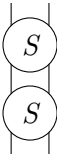
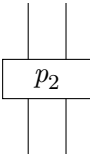
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▶  $= \delta \cdot$ 

▶  $= \text{zero}$

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▶  $= i \cdot$ 

▶  $=$ 

Summary

We have just proved the following:

Theorem

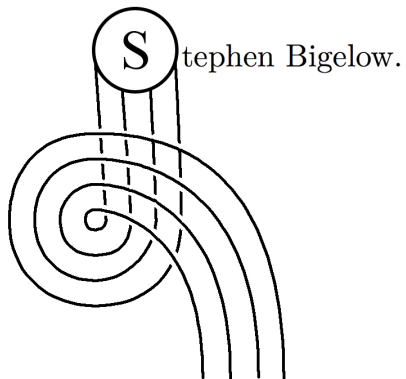
The defined planar algebra is not trivial

This is part of the **Kuperberg program**:

Give a presentation for every interesting planar algebra, and prove as much as possible about the planar algebra using only its presentation.

Thank You

I also want to thank my advisor



$$[m+1] \cdot \begin{array}{c} \vdots \\ \vdots \\ \boxed{p_{m+1}} \\ \vdots \\ \vdots \end{array} = [m+1] \cdot \begin{array}{c} \vdots \\ \vdots \\ \boxed{p_m} \\ \vdots \\ \vdots \end{array} \Big| - [m] \cdot \begin{array}{c} \vdots \\ \vdots \\ \boxed{p_m} \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \boxed{p_m} \\ \vdots \\ \vdots \end{array}$$

$$[m] \cdot \begin{array}{c} \vdots \\ \vdots \\ \boxed{p_m} \\ \vdots \\ \vdots \end{array} \Big| = [m+1] \cdot \begin{array}{c} \vdots \\ \vdots \\ \boxed{p_{m-1}} \\ \vdots \\ \vdots \end{array}$$

